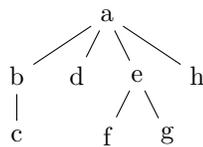


# Foundations of XML Types

## Trees and Tree Automata

1. Every tree can be represented as a binary tree. Give the binary tree associated with the tree shown below by the encoding of trees into binary trees used during the course:



2. Give a bottom-up deterministic tree automaton that recognize the tree language  $L$  composed of the two trees below:



3. Bottom-up tree automata seen during the course traverse trees from the leaves to the root. In a similar manner, one may define top-down tree automata that recognize trees by going in the opposite direction: from the root to the leaves. Specifically, a top-down tree automaton  $A$  consists in:

$\text{Alphabet}(A)$ : finite alphabet of symbols  
 $\text{States}(A)$ : finite set of states  
 $\text{Rules}(A)$ : finite set of transition rules  
 $\text{Initial}(A)$ : finite set of initial states ( $\subseteq \text{States}(A)$ )  
 $q_{\text{acc}} \in \text{States}(A)$ : final state

There are two major differences with automata seen during the course:

- transition rules are either of the form:  $q \xrightarrow{a} (q_1, q_2)$  where  $q, q_1, q_2 \in \text{States}(A)$  and  $a \in \text{Alphabet}(A)$  or of the form  $q \xrightarrow{\epsilon} q_1$  for leaves.
- a tree is accepted if and only if there exists a run for which all the leaves are labeled with  $q_{\text{acc}}$ .

Give a top-down tree automaton that recognizes  $L$ .

4. Do you see any interest of top-down tree automata in the context of XML stream processing where XML documents are sequentially parsed (only once) and processed on the fly? Explain.
5. A top-down tree automaton is deterministic iff (1) there is at most one initial state and (2) for each  $q \in \text{States}(A)$  et  $a \in \text{Alphabet}(A)$  there is at most one rule  $q \xrightarrow{a} (q_1, q_2)$  (intuitively, there is at most one possible transition for each state and symbol).  
Is it possible to give a deterministic top-down tree automaton that recognize  $L$ ? Either give it or justify.

6. It is known that non-deterministic bottom-up and non-deterministic top-down automata are equally expressive. From your answers to the previous questions, what can you conclude about the respective expressive power of deterministic bottom-up and deterministic top-down tree automata? Justify.